

*Thin Objects*  
*An Abstractionist Account*  
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# Abstractionism

Two main strands:

**(M)** Mathematics – or substantial amounts of it – may be based on **abstraction principles** (APs), or similar:

$$\forall\alpha\forall\beta(\S(\alpha) = \S(\beta) \leftrightarrow \alpha \sim \beta) \quad (\text{AP})$$

$\alpha, \beta$  – ‘specifications’     $\alpha \sim \beta$  – ‘unity relation’     $\S(\alpha)$  – ‘abstract’

**(P)** APs – or ‘good’ APs – enjoy a **privileged epistemic status**

# Frege: the abstractionist prototype

( $\mathbf{M}_F$ ) BLV + SOL interprets PA  
 $\text{ext } F = \text{ext } G$  iff  $F$  and  $G$  are coextensive (BLV)

SOL – full second-order logic

( $\mathbf{P}_F$ ) BLV is a law of logic

But we all know how that went:

**Bad company #1:** BLV + SOL  $\vdash \perp$  (Russell's paradox)

– assuming SOL is okay (as we shall henceforth):

- $\mathbf{M}_F$  is trivial
- $\mathbf{P}_F$  is untenable

# Hale and Wright: the neo-Fregean programme

(**M<sub>HW</sub>-1**) HP + SOL interprets PA (Frege's theorem)  
 $\#F = \#G$  iff  $F$  and  $G$  are equinumerous (HP)

(**M<sub>HW</sub>-2**) New V + SOL interprets a substantial subtheory of ZF  
 $\text{ext } F = \text{ext } G$  iff  $F$  and  $G$  are both universe-sized or coextensive (New V)

(**P<sub>HW</sub>**) HP and other good APs are **'implicit definitions'**

**Bad company #2:** which APs are good?

	Satisfiable		
Scylla ↑	Unbounded	Conservative	↑ inconsistent
Charybdis ↓	Stable	Irenic	↓ can't interpret ZFU without relativization
	Strongly Stable		

## A different approach – predicative vs impredicative APs

**Neo-Fs:** APs may have an ‘**impredicative**’ character – e.g.:

$$\{xx\} = \{yy\} \leftrightarrow \forall z(z < xx \leftrightarrow z < yy) \quad (\text{V})$$

–  $\{xx\}$  and  $\{yy\}$  must denote something in domain of  $\forall z$

neo-Fs defend impredicative APs via  $\forall$ -absolutism:

**$\forall$ -absolutism:**  $\forall$  ranges over an absolutely comprehensive domain

**Linnebo:** APs should be ‘**predicative**’ – e.g.

$$\{xx\} = \{yy\} \leftrightarrow \forall z(z < xx \leftrightarrow z < yy) \quad (2\text{V})$$

–  $\{xx\}$  and  $\{yy\}$  may denote something outside domain of  $\forall z$

# Linnebo's abstractionism

(P<sub>L</sub>) 2V and other 2APs are ‘“free truths”’

$$\forall\alpha\forall\beta(\S(\alpha) = \S(\beta) \leftrightarrow \alpha \sim \beta) \quad (2AP)$$

“free truths”? – community C lay down assertibility-conditions, e.g.:

(AC)  $\forall\alpha\forall\beta: \ulcorner \S(\alpha) = \S(\beta) \urcorner$  is assertible of  $\alpha$  and  $\beta$  iff  $\alpha \sim \beta$

ØL argues – best interpretation of C renders 2AP true and knowable

– let's just **grant P<sub>L</sub>**:

PFO – two-sorted plural logic

**Bad company #3** – a ‘simple and definitive’ solution?

- 2V + PFO  $\not\vdash \perp$  – ‘just about any’ 2AP okay
- 2V + PFO  $\vdash \exists y\forall x(x \neq y)$  –  **$\forall$ -absolutism fails**
- 2AP + PFO  $\not\vdash$  “there are more than two abstracts”

**Q1:** how do we overcome weakness of 2APs to obtain PA or ZF?

# Q1: overcoming weakness of predicative abstraction (ZF)

**A1 (ZF): iterating** 2V-style abstraction ( $\infty$ -ly many times)

$\square \phi$ : ‘however we abstract,  $\phi$ ’     $\diamond \phi$ : ‘we can abstract so that  $\phi$ ’

**MS:** MPFO + Foundation + six more axioms:     $\square$  – interpretational

A1.  $\square \forall uu \diamond \exists x \text{SET}(uu, x)$

A2.  $x = y \leftrightarrow \forall u(u \in x \leftrightarrow u \in y)$

A3.  $\square \forall y \exists xx \square \forall x(x < xx \leftrightarrow x \in y)$

A4.  $\square \forall y \exists xx \square \forall x(x < xx \leftrightarrow x \subseteq y)$

A5.  $\square \forall \vec{v}(\phi^\diamond(\vec{v}) \rightarrow \diamond \phi(\vec{v}))$

A6.  $\text{fn}[\phi]^\diamond \rightarrow \square \forall xx \diamond \exists yy(\forall x < xx)(\exists y < yy)(\phi^\diamond(x, y))$

**(M<sub>L</sub>-1)** MS interprets ZF

– and, although  **$\forall$ -absolutism** fails,  $\emptyset L$  endorses:

**$\square \forall$ -absolutism:**  $\square \forall$  and  $\diamond \exists$  generalize about the whole hierarchy

**Q2:** even if 2APs are, are MS-axioms “free truths” (or similar)?

# Q1: overcoming weakness of predicative abstraction (PA)

**A1 (PA):** one round of **modal** predicative abstraction

Two key assumptions (for suitable  $\S$  and  $\sim$ ):  $\Box$  – metaphysical

$$\Box \forall \alpha \forall \beta (\S(\alpha) = \S(\beta) \leftrightarrow \alpha \sim \beta) \quad (\Box 2AP)$$

$$\Box \forall x (\text{ABST}_{\S}(x) \rightarrow \Box \exists y (y = x)) \quad (\Box E)$$

**ØL:**  $\Box E$  – ‘very plausible’ and ‘very widely shared’

– but  $\Box E$  conflicts with another plausible assumption:

**No Specificationless Abstracts:** an abstract item exists only if some of its specifications exist – e.g.  $\{\text{ØL}\}$  exists only if  $\text{ØL}$  exists

**Q3:** how can we ‘introduce’ an abstract without its specification?



### Q3: ‘introducing’ specificationless abstracts?

**A3 [suggested]:** predicative abstraction on **possible specifications**

Frame assertibility-conditions with a modal language (with @):

**(AC<sup>★</sup>)**  $\Pi\alpha\Pi\beta$ : @  $\ulcorner \S(\alpha) = \S(\beta) \urcorner$  is assertible of  $\alpha$  and  $\beta$  iff  $\alpha \sim \beta$

$$\Pi\alpha = @\Box\forall\alpha@$$

– adapting ØL’s style of argument...

**(P<sub>L</sub><sup>★</sup>)** 2AP<sup>★</sup>s are **“free truths”**

$$\Pi\alpha\Pi\beta(@(\S(\alpha) = \S(\beta)) \leftrightarrow \alpha \sim \beta) \quad (2AP^{\star})$$

– cost: renounce part of ØL’s view?

□ – interpretational

**□∀-absolutism fails:**  $2AP^{\star} + MPFO_{@} \vdash @\exists yy\Box\forall xx(xx \neq yy)$

## 2APs – a ‘simple and definitive’ solution to bad company?

We’ve been granting:

**(P<sub>L</sub>)** 2APs are “free truths”

Questions:

**Q1:** how do we overcome weakness of 2APs to obtain PA or ZF?

**A1 (ZFC): transfinite iteration** of predicative abstraction

– **Q2:** are the MS-axioms “free truths”?

**A1 (PA): modal** predicative abstraction

– **Q3:** how do we introduce specificationless abstracts?

– **A3 [suggested]:** abstraction on **possible specifications**

– **Q4:** are 2AP\*s “free truths” (contrary to  $\Box\forall$ -absolutism)?

## Appendix: MPFO@

The system MPFO@ comprises the following:

- a free, two-sorted formulation of PFO
- a normal modal system for  $\Box$  and  $@$
- further axioms governing  $@$ :
  - a1:  $@\neg\phi \leftrightarrow \neg@ \phi$
  - a2:  $@(\phi \rightarrow \psi) \rightarrow (@\phi \rightarrow @\psi)$
  - a3:  $@(@\phi \rightarrow \phi)$
  - a4:  $\Diamond @\phi \rightarrow @\phi$
  - a5:  $@[\forall v @\phi \leftrightarrow @\forall v \phi]$